On the Benefits of Randomly Adjusting Anytime Weighted A*

Abstract

- Anytime Weighted A* (Hansen and Zhou, 2007; Hansen, Zilberstein, and Danilchenko, 1997) is an anytime heuristic search algorithm that uses a weight to scale the heuristic to manage the trade-off between solution quality and running time.
- We propose a randomized version of this algorithm, called Randomized Weighted A^* , that randomly adjusts its weight at runtime
- RWA* typically outperforms AWA* with static weights on a range of benchmark problems.

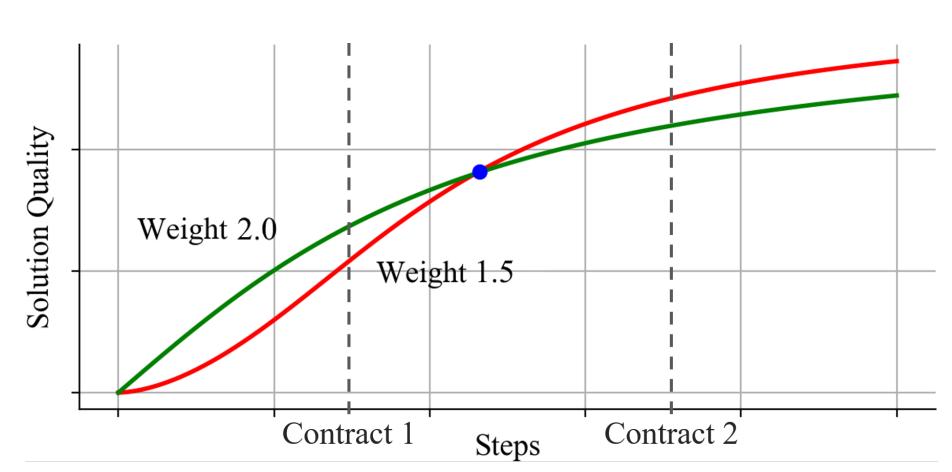


Figure: An example of two executions of AWA*. Weight=2.0 is better when the available running time is contract 1.

Introduction

- **Contract setting**: A fixed computation time is available to solve a problem.
- **Trade-off**: Higher weights lead to better solutions in short-term.
- **Best weight** depends on the characteristics of the domain, the details of the instance, and the available computation time (contract duration).
- Tune best static weight for a problem (Hansen and Zhou, 2007)
- Tune at runtime heuristically (Sun, Druzdzel, and Yuan, 2007; Thayer and Ruml, 2009, 2008)
- Adjust at runtime using deep-RL (Bhatia, Svegliato, and Zilberstein, 2021)
- Adjust at runtime randomly? Advantage: simplicity, no hyperparameters, no offline experimentation.

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Randomized Weighted A*

- For each node expansion, RWA* samples a weight uniformly from a fixed set of weights e.g., $w \sim W = \{1, 1.5, 2, 3, 4, 5\}.$
- RWA* maintains an open list corresponding to each w. Same nodes, different priorities.
- Operate on the open lists in parallel for efficiency.

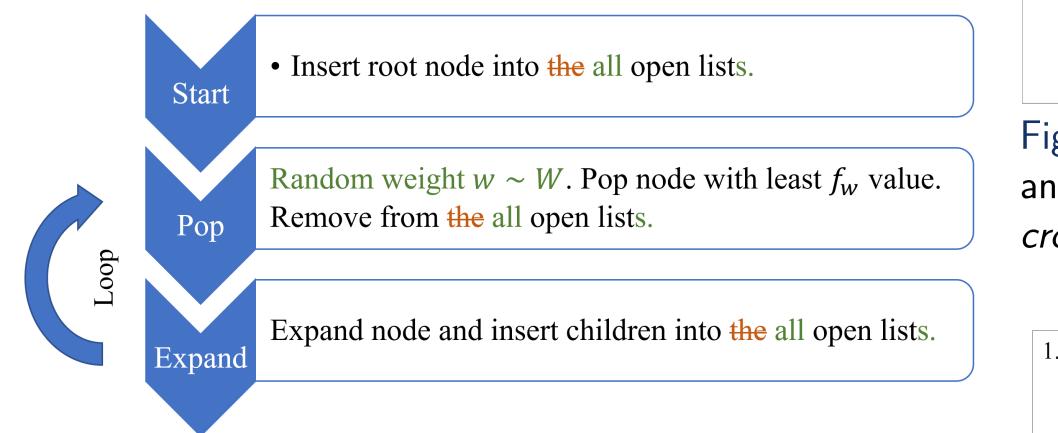


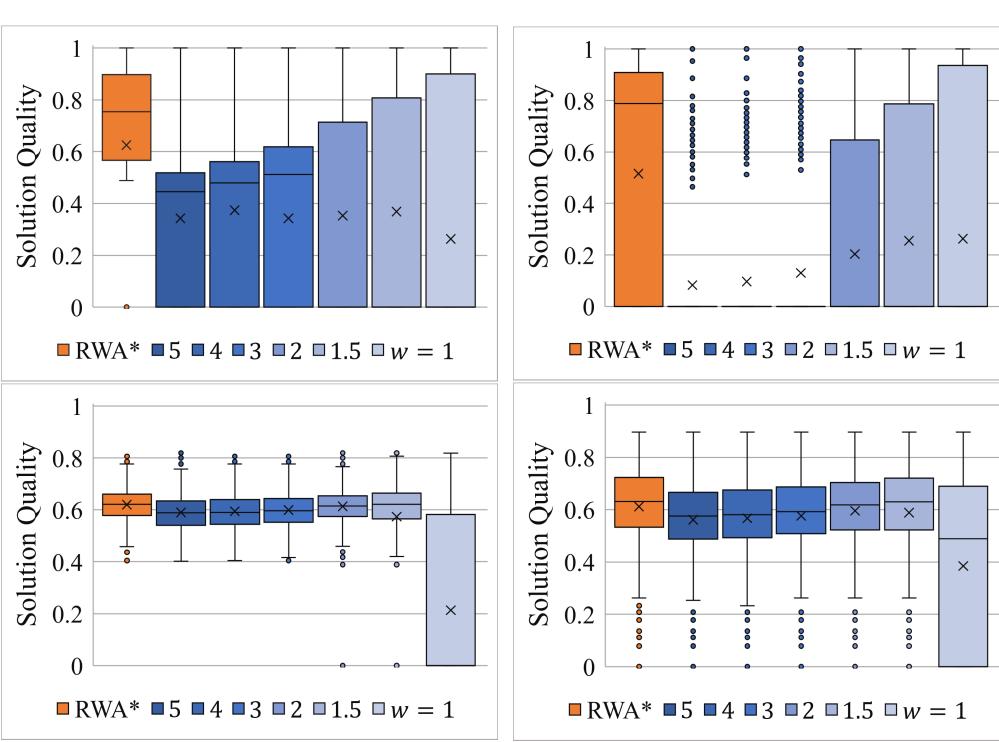
Figure: Differences between RWA* and AWA*

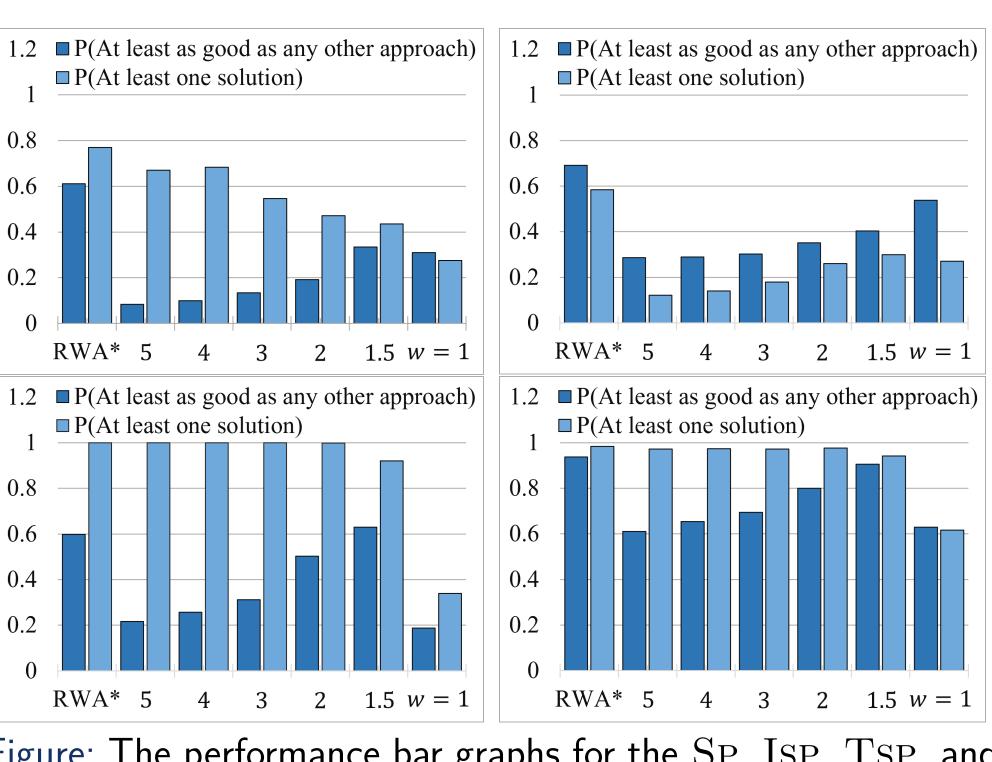
Experimental Setup

- RWA* $\mathbf{w} \sim \{1, 1.5, 2, 3, 4, 5\}$ vs AWA* with static weights 1, 1.5, 2, 3, 4, 5 (commonly used).
- **Domains**: Sliding-Puzzle (SP), Inverse-Sliding-Puzzle (ISP), Travelling-Salesman-Problem (TSP), City-Navigation-Problem (CNP)
- **500 instances** per domain of varying difficulty.
- Node-expansions budget (contract) of 6000 for SP, ISP and 3000, 2400 for TSP, CNP.

Results

- RWA* computes solutions with a **higher** quality on average than any static weight,
- 2 has the highest probability of computing a solution that is at least as good as any other approach,
- (3) has the highest probability of computing at least one solution compared to any static weight.





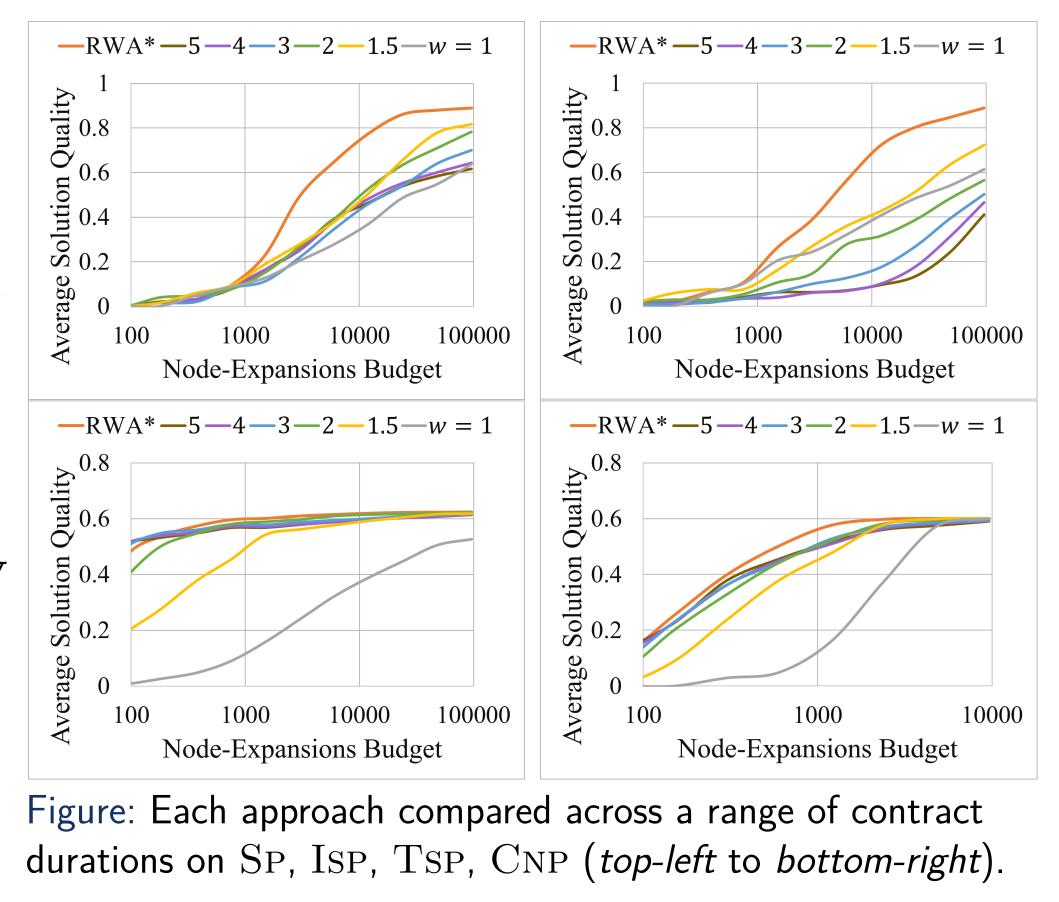
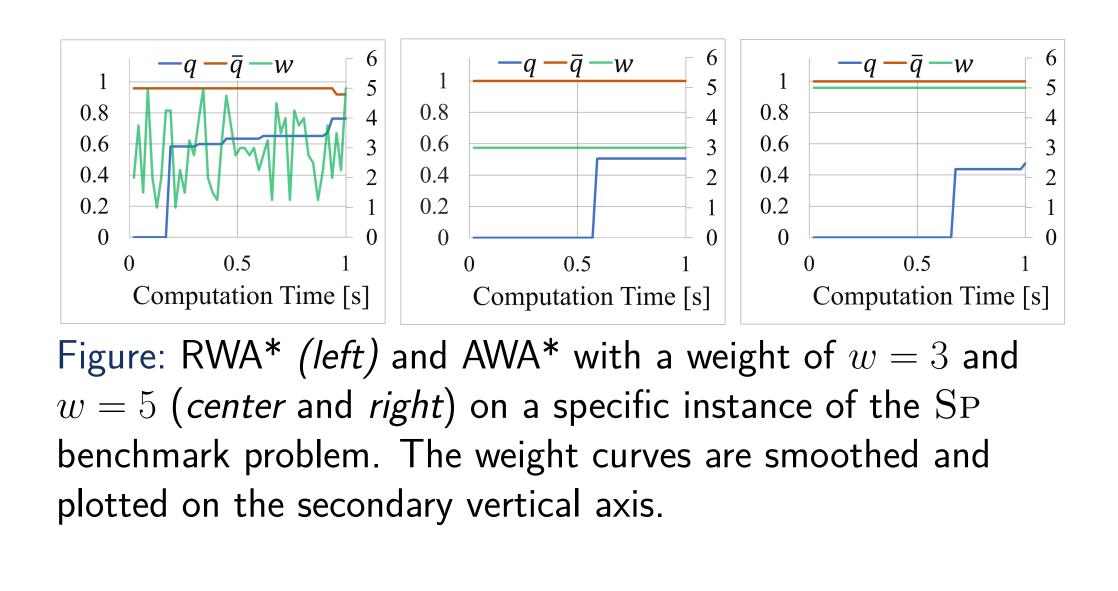


Figure: The solution quality box plots for the SP, ISP, TSP, and CNP benchmark problems (*top-left* to *bottom-right*). The crosses denote the mean and the *bullets* denote the outliers.





Conclusion

• RWA* (i) computes better solutions on average, (ii) exhibits a higher probability of computing any solution at all, and (iii) exhibits a higher probability of computing a solution at least as good as any static weight of AWA* in a contract setting across a range of contract durations on our benchmark domains.

• RWA* is appealing because it is easy to implement and effective without any extensive offline experimentation or parameter tuning.

References

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